

MORE ON SIMPLICIAL CATS

Recall from last week:

- **Monoidal categories**, functors, nat. transformations, etc.
- **Enriched categories**, functors, nat. transformations, etc.

Lemma. We have a 2-functor

$$\text{MonCat} \longrightarrow \text{Cat}$$

$$\Phi: V \rightarrow V' \rightsquigarrow \Phi_*: \text{Cat}_V \rightarrow \text{Cat}_{V'} \text{ a 1-functor}$$

$$\text{ob } \Phi_* \mathcal{C} = \text{ob } \mathcal{C}, \quad \text{Hom}_{\Phi_* \mathcal{C}}(x, y) = \Phi \text{Hom}_{\mathcal{C}}(x, y).$$

This sends monoidal adjunctions to adjunctions.

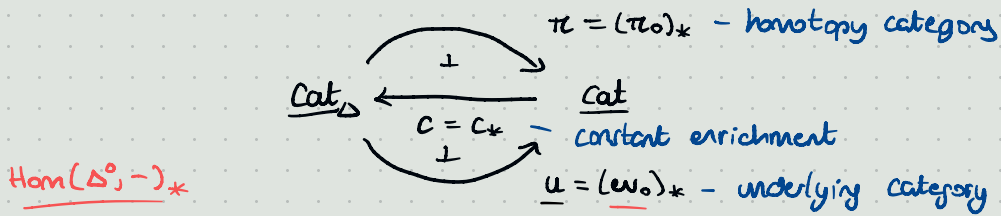
Def: Cat_{Δ} is the category of simplicial categories, i.e. categories enriched over $s\text{Set}$.

Lemma: Cat_{Δ} is bicomplete.

Recall from earlier: we have adjunctions

$$\begin{array}{ccc} & \begin{array}{c} \xrightarrow{\pi_0} \\ \perp \\ \xleftarrow{c} \\ \perp \\ \xrightarrow{ev_0} \end{array} & \\ s\text{Set} & & \text{Set} \\ & \begin{array}{c} \text{connected components} \\ \text{constant simplicial set} \\ \text{evaluation at } [0] \end{array} & \end{array}$$

Lemma. These all of these functors are monoidal, and so are the adjunctions, so we get adjunctions



Idea of proof: both $\underline{\text{Set}}$ and $\underline{\text{Set}}$ are cartesian ($\otimes = \times$), so we have canonical maps

$$\begin{array}{ccc}
 X \times Y & \begin{array}{l} \rightarrow X \\ \rightarrow Y \end{array} & \rightsquigarrow F(X \times Y) \begin{array}{l} \rightarrow FX \\ \rightarrow FY \end{array} \rightsquigarrow F(X \times Y) \rightarrow FX \times FY.
 \end{array}$$

- C and ev_0 are right adjoints, so they preserve products.
- For π_0 , consider

$$\begin{array}{ccc}
 (X \times Y)_0 & \xrightarrow{\cong} & X_0 \times Y_0 \\
 \downarrow & & \downarrow \\
 \pi_0(X \times Y) & \longrightarrow & \pi_0 X \times \pi_0 Y
 \end{array}$$

↑ must be surjective
↖ also injective since we can lift relations, because ev_0 is monoidal. \square

Terminology. Let $\mathcal{C} \in \underline{\text{Cat}}_{\Delta}$, a morphism $x \rightarrow y$ in \mathcal{C} is a morphism $x \rightarrow y$ in $u\mathcal{C}$, i.e. a 0-simplex of $\text{Hom}_{\mathcal{C}}(x, y)$. A morphism in \mathcal{C} is an equivalence if its image in $\pi\mathcal{C}$ is an isomorphism.

Next: defining the nerve of a simplicial category.

We do this by constructing a cosimplicial object in Cat_Δ .

$$\Delta \rightarrow \text{Cat}_\Delta$$

Def. Let J be a finite non-empty totally ordered set, $i, j \in J$.

$$\begin{aligned} P_{i,j} &= \{ I \subseteq J : i, j \in I \text{ and } k \in I \Rightarrow i \leq k \leq j \} \\ &= \{ \text{subsets of } [i, j] \text{ containing the endpoints} \}. \end{aligned}$$

$P_{i,j}$ is a poset under inclusion.

If $i > j$, then $P_{i,j} = \emptyset$. If $i \leq j \leq k \in J$, then

$$N(P_{i,j}) \times N(P_{j,k}) \rightarrow N(P_{i,k})$$

$$(I, I') \mapsto I \cup I'$$

is an associative binary operation.

[n]

For each J , we have a simplicial category $\mathcal{C}[\Delta^J]$:

$$\text{ob } \mathcal{C}[\Delta^J] = J, \quad \text{Hom}_{\mathcal{C}[\Delta^J]}(i, j) = \begin{cases} \emptyset & \text{if } i > j \\ N(P_{i,j}) & \text{o/w.} \end{cases}$$

Composition is induced by the binary operation above.

Lemma. Let $n \geq 1$, $J = [n]$. Then

(a) $N(P_{0,n}) \cong (\Delta^1)^{n-1}$

(b) $P_{i,j} \cong P_{0, j-i}$.

Proof. (b) is clear. For (a), sufficient to show

$$P_{0,n} \cong \underbrace{[1] \times \cdots \times [1]}_{n-1 \text{ times}}$$

↑

$$s \mapsto (x_s^{(1)}, \dots, x_s^{(n-1)})$$

← an element of this represents a subset of $(0, n)$ in $[n]$.

Lemma. If \mathcal{C} is a category with an initial or terminal object, then $N(\mathcal{C})$ is contractible. $\emptyset \in \mathcal{C} \text{ initial} \Rightarrow \emptyset \Rightarrow 1_{\mathcal{C}}$
 $\mathcal{C} \times [1] \rightarrow \mathcal{C} \xrightarrow{\sim} N(\mathcal{C}) \times \Delta^1 \rightarrow N(\mathcal{C})$

Cor. For $i \leq j$, $\text{Hom}_{\mathcal{C}[\Delta^j]}(i, j)$ is contractible.

Lemma. There is a unique isomorphism $\pi: \mathcal{C}[\Delta^n] \cong [n]$ which is the identity on objects. Hence, we have a canonical functor $\mathcal{C}[\Delta^n] \rightarrow \mathcal{C}[n]$. $\mathcal{C}[\Delta^n]$

Lemma. The assignment $J \mapsto \mathcal{C}[\Delta^J]$ defines a functor $\text{Linordset} \rightarrow \text{Cat}_{\Delta}$.

In particular, we have a cosimplicial object $\Delta \rightarrow \text{Cat}_{\Delta}$ given by $[n] \rightarrow \mathcal{C}[\Delta^n]$.

Idea of proof: Unfolding definitions, it suffices to construct for any monotone map $f: J \rightarrow J'$ and $i \leq j \in J$, a monotone map $P_{i,j} \rightarrow P_{f(i), f(j)}$. This is $I \mapsto f(I)$. \square

Def. For $\mathcal{C} \in \text{Cat}_{\Delta}$, its **simplicial nerve** (or homotopy-coherent nerve) is the sset:

$$N(\mathcal{C})_n = \text{Hom}_{\text{Cat}_{\Delta}}(\mathcal{C}[\Delta^n], \mathcal{C})$$

Lemma. If $\mathcal{C} \in \text{Cat}$, then $N(\mathcal{C}) \cong N(c\mathcal{C})$.

Proof: $N(c\mathcal{C})_n = \text{Hom}_{\text{Cat}_{\Delta}}(\mathcal{C}[\Delta^n], c\mathcal{C})$
 $\cong \text{Hom}_{\text{Cat}}(\pi \mathcal{C}[\Delta^n], \mathcal{C})$
 $\cong \text{Hom}_{\text{Cat}}([n], \mathcal{C}) = N(\mathcal{C})_n$.

Discussion. What do (low dimensional) simplices in $N(\mathcal{C})$ look like?

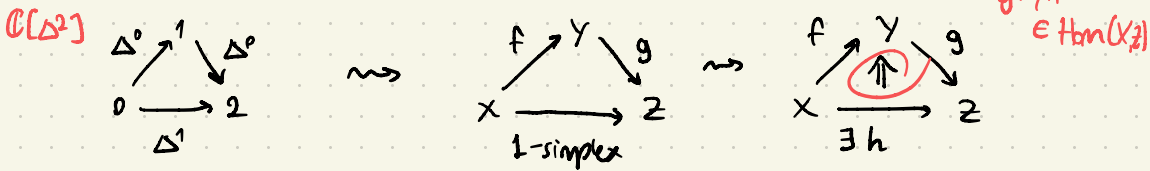
- $\mathcal{C}[\Delta^0]$ has a single object and $\text{Hom}_{\mathcal{C}[\Delta^0]}(0,0) = N(P_{0,0}) = \Delta^0$. 0,0
u
- $\mathcal{C}[\Delta^1]$ has two objects and all hom-sets are Δ^0 . P_{0,1}

Hence, $\mathcal{C}[\Delta^0] = \mathcal{C}[0]$ and $\mathcal{C}[\Delta^1] = \mathcal{C}[1]$.

- $N(\mathcal{C})_0 = \text{Hom}_{\text{Cat}_{\Delta}}(\mathcal{C}[\Delta^0], \mathcal{C}) \cong \text{Hom}_{\text{Cat}}([0], \mathcal{U}\mathcal{C}) \cong \text{ob } \mathcal{U}\mathcal{C} = \text{ob } \mathcal{C}$.
- $N(\mathcal{C})_1 = \text{Hom}_{\text{Cat}_{\Delta}}(\mathcal{C}[\Delta^1], \mathcal{C}) \cong \text{Hom}_{\text{Cat}}([1], \mathcal{U}\mathcal{C}) \cong \text{mor } \mathcal{U}\mathcal{C} = \text{mor } \mathcal{C}$.

What is $\mathcal{C}[\Delta^2]$? It has three objects. All hom-sets are Δ^0 except $\text{Hom}_{\mathcal{C}[\Delta^2]}(0,2) = N(P_{0,2}) = \Delta^1$. {0,2}, {0,1,2}

Hence, a simplicial functor $\mathcal{C}[\Delta^2] \rightarrow \mathcal{C}$ picks out:



Lemma: There is a unique colimit preserving functor

$$\mathcal{C}[-] : \text{sSet} \rightarrow \text{Cat}_{\Delta}$$

which sends $\Delta^n \mapsto \mathcal{C}[\Delta^n]$. It is left-adjoint to N .

Fact: $\mathcal{C}[-]$ preserves monomorphisms, i.e. if $A \subseteq X$ is an inclusion of ssets then $\mathcal{C}[A] \hookrightarrow \mathcal{C}[X]$.

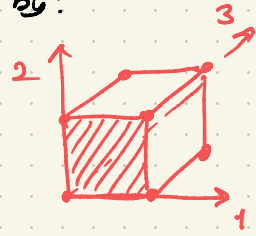
Lemma. Let $0 < j < n$. $\mathcal{C}[\Delta_j^n] \subseteq \mathcal{C}[\Delta^n]$ is given by:

(1) $\text{ob } \mathcal{C}[\Delta_j^n] = \text{ob } \mathcal{C}[\Delta^n]$.

(2) $\text{Hom}_{\mathcal{C}[\Delta_j^n]}(i, k) = \text{Hom}_{\mathcal{C}[\Delta^n]}(i, k)$ except

$\text{Hom}_{\mathcal{C}[\Delta_j^n]}(0, n) \neq \text{Hom}_{\mathcal{C}[\Delta^n]}(0, n)$

which is given by the subsimplicial set of $(\Delta^1)^{n-1}$ obtained by deleting the interior of the bottom j -face.



Proof idea: $\Delta_j^n = \bigcup_{i \neq j} \Delta^{n-1} i i s$ so look at $\mathcal{C}[\Delta^{n-1} i i s] \subseteq \mathcal{C}[\Delta^n]$.

Lemma. $N(\mathcal{C})$ is a complex. Further, if all n -sets in \mathcal{C} are Kan complexes then $N(\mathcal{C})$ is an ∞ -category.

Proof:

$$\begin{array}{ccc} I^n & \rightarrow & N(\mathcal{C}) \\ \downarrow & \dashrightarrow & \\ \Delta^n & & \end{array} \quad \Leftrightarrow \quad \begin{array}{ccc} \mathcal{C}[I^n] & \rightarrow & \mathcal{C} \\ \downarrow & \dashrightarrow & \\ \mathcal{C}[\Delta^n] & & \end{array}$$

But we have a retraction of $\mathcal{C}[I^n] \rightarrow \mathcal{C}[\Delta^n]$.

To see this, note $\mathcal{C}[I^n] = \mathcal{C}[n]$ since $I^n = I^{n-1} \cup_{\Delta^0} I^1$.

But we saw there is a unique simplicial functor

$\mathcal{C}[\Delta^n] \rightarrow \mathcal{C}[n]$ that is the identity on objects. Then

$$\mathcal{C}[n] \cong \mathcal{C}[I^n] \hookrightarrow \mathcal{C}[\Delta^n] \rightarrow \mathcal{C}[n]$$

is the identity.

For Kan complexes, need $0 < j < n$

$$\begin{array}{ccc}
 \Delta_j^n & \longrightarrow & N(\mathcal{C}) \\
 \downarrow & \nearrow & \longleftarrow \\
 \Delta^n & \dashrightarrow & \mathcal{C}[\Delta_j^n] \xrightarrow{f} \mathcal{C} \\
 & & \downarrow \nearrow \\
 & & \mathcal{C}[\Delta^n]
 \end{array}$$

From the previous lemma, we only need to worry about

$$\begin{array}{ccc}
 \text{Hom}_{\mathcal{C}[\Delta_j^n]}(0, n) & \xrightarrow{f} & \text{Hom}_{\mathcal{C}}(f(0), f(n)) \\
 \downarrow & \nearrow & \\
 \text{Hom}_{\mathcal{C}[\Delta^n]}(0, n) & \dashrightarrow &
 \end{array}$$

But this follows from the fact that the vertical map is **anodyne**. One checks that this extension in fact defines a simplicial functor. □

Remark. Nerves of categories enriched over Kan complexes will be $(\infty, 1)$ -categories?

Def. Consider the simplicial category of CW-complexes and hom-sets the singular set of the mapping space. Its simplicial nerve is the **∞ -category of spaces, \mathcal{Spc}** .

Lemma. The product and coproduct of ∞ -cats is an ∞ -cat.

Proof. For product, find extension in each factor and then put together. For coproduct, note that Δ_j^n and Δ^n are connected. □

Def. A **sub- ∞ -cat** $\mathcal{C}' \subseteq \mathcal{C}$ is a subsimplicial set determined by subsets $X \subseteq \mathcal{C}_0$ and $S \subseteq \mathcal{C}_1$, between objects in X and closed under composition and equivalences. An n -simplex of \mathcal{C} belongs in \mathcal{C}' iff the edges of its restriction to the spine I^n are in S . A subcategory is **full** if S contains all 1-simplices whose boundary is in X .

Lemma. A sub- ∞ -category is an ∞ -cat. Its homotopy category is the subcategory of $h\mathcal{C}$ on the image of S .

$$\begin{array}{ccc} \mathcal{C}' & \longrightarrow & \mathcal{C} \\ \downarrow & & \downarrow \\ N(h\mathcal{C}') & \longrightarrow & N(h\mathcal{C}) \end{array}$$

is a pullback. For any subcategory $\mathcal{D} \subseteq h\mathcal{C}$, this pullback gives a sub- ∞ -cat.

Cor. There is a 1-to-1 correspondence between sub- ∞ -cats of \mathcal{C} and subcategories of $h\mathcal{C}$. Full correspond to full.

Def. A **natural transformation** between functors $f, g: \mathcal{C} \rightarrow \mathcal{D}$ is a simplicial map $\mathcal{C} \times \Delta^1 \rightarrow \mathcal{D}$ that restricts to f and g .

Remark. This is exactly a 1-simplex in $\text{Hom}(\mathcal{C}, \mathcal{D})$.

We will see that $\text{Hom}(\mathcal{C}, \mathcal{D})$ is also an ∞ -cat and $N(\text{Fun}(\mathcal{C}, \mathcal{D})) = \text{Hom}(N\mathcal{C}, N\mathcal{D})$ since N is fully faithful.

$$- \times A \rightarrow \text{Hom}(A, -)$$

§ Analyse d. (kan.) Fibrationen $C = \text{sSets}$

Def: A (left/right/inner) fibration $p: K \rightarrow S$ has the

RLP wrt \mathbb{A} left/right/inner horns.

right
lifting



$\text{Sclor}(C)$

$$\mathcal{X}_R(S) = \{ f: X \rightarrow C \mid f \text{ has RLP wrt } S \} \quad \mathcal{X}_L(S) \text{ (LRP)}$$

$$\mathcal{X}_R(\{\text{left horns}\}) = \text{Left Fibration.}$$

Def: A map $f: A \rightarrow B$ is (left/right/inner) anodyne if it has LLP wrt (\dots) fibrations. i.e. $\mathcal{X}_L(\mathcal{X}_R(\text{left/right/inner horns}))$

$$S \rightarrow S \subset \mathcal{X}_L(\mathcal{X}_R(S)) =: \mathcal{X}(S)$$

Def: A $\text{Sclor}(C)$ is saturated if:

1) it contains all isos

2) Stable under pushouts:

$$\begin{array}{ccc} A & \rightarrow & A' \\ S \ni i \downarrow & & \downarrow i' \\ B & \rightarrow & B' \end{array} \Rightarrow i' \in S$$

3) Stable retracts:

$$\begin{array}{ccccc} & \xrightarrow{\text{id}} & & & \\ A' & \rightarrow & A & \rightarrow & A' \\ \downarrow & \parallel & \downarrow i & \parallel & \downarrow i' \\ B' & \rightarrow & B & \rightarrow & B' \\ & \xrightarrow{\quad} & & & \end{array} \quad \text{If } i \in S, \text{ so is } i'.$$

4) Stable under composable composition (N-diagram)

$$\begin{array}{ccccccc} A_0 & \xrightarrow{i_0} & A_1 & \xrightarrow{i_1} & A_2 & \rightarrow & \dots \\ & \searrow i_0 & \downarrow i_1 & \swarrow i_2 & & & \\ & \Rightarrow \in S & \text{colim } A_i & & & & \end{array} \quad i_j \in S \quad \forall j.$$

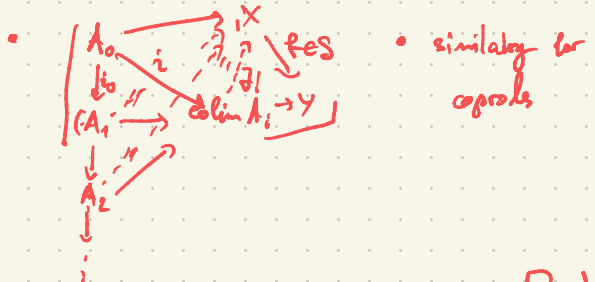
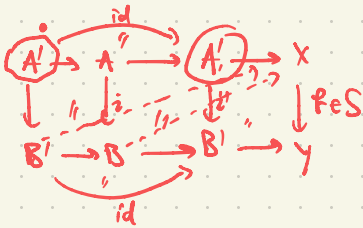
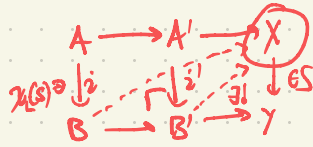
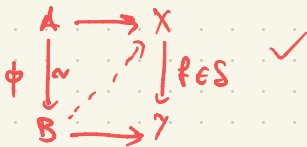
5) Coproduct: If $\{A_i \xrightarrow{i_i} B_i\}_{i \in J} \in \mathcal{S}$

$$\Rightarrow \coprod_i A_i \xrightarrow{\coprod i_i} \coprod_i B_i \in \mathcal{S}$$

Def $\mathcal{S} \subseteq \text{Mor}(\mathcal{C})$. The saturated closure of \mathcal{S} is the smallest saturated set containing \mathcal{S} , i.e. $\bar{\mathcal{S}} = \bigcap_{\substack{T \text{ saturated} \\ T \supseteq \mathcal{S}}} T$

Lemma $\mathcal{S} \subseteq \text{Mor}(\mathcal{C})$. Then, $\mathcal{R}_L(\mathcal{S})$ is saturated.

Prf:



Prop: (Small Object Argument) $\mathcal{C} = \text{sSets}$

□

Let $\mathcal{S} = \{A_i \rightarrow B_i\}_{i \in J} \subseteq \text{Mor}(\mathcal{C})$ s.t. A_i 's have finitely many non-deg. simplices

Then any $F: X \rightarrow Y$ in sSets can be factorized.

$$X \xrightarrow{F} Y$$

$$\wr \downarrow \nearrow \wr$$

$$h \downarrow \quad g$$

s.t. $h \in \bar{\mathcal{S}}, g \in \mathcal{R}_R(\mathcal{S})$.